Key

n = the original input size T(n) = time required to process input of size n F = factor by which the input size is multiplied T(Fn) = time required to process input of size Fn

O(log n) Algorithms

- It doesn't matter which base you use for the logs, as long as you use the same base throughout the problem.
- If the numbers in the problem are powers of 10, use base 10.
- If the numbers in the problem are powers of 2, use base 2.
- Even if you don't use the base suggested above, you'll still get the same answer, but the arithmetic will be more complicated.

 $T(n) = c \log(n)$ $T(Fn) = c \log(Fn)$

T(Fn)	_(1)	$\log(F)$
T(n)	-(1+	$\log(n)$

 $\frac{T(n)}{\log(n)} = \frac{T(Fn)}{\log(Fn)}$

O(n) Algorithms

T(n) = cnT(Fn) = c(Fn)

 $\frac{T(Fn)}{T(n)} = F$

 $\frac{T(n)}{n} = \frac{T(Fn)}{Fn}$

O(n log n) Algorithms

 $T(n) = c n \log(n)$ T(Fn) = c (Fn) log(Fn)

$$\frac{T(Fn)}{T(n)} = F\left(1 + \frac{\log(F)}{\log(n)}\right)$$
$$\frac{T(n)}{n\log(n)} = \frac{T(Fn)}{(Fn)\log(Fn)}$$

O(n²) Algorithms

$$T(n) = c n2$$

T(Fn) = c (Fn)²

$$\frac{T(Fn)}{T(n)} = F^2$$

$$\frac{T(n)}{n^2} = \frac{T(Fn)}{(Fn)^2}$$

O(n³) Algorithms

T(n) = c n³T(Fn) = c (Fn)³

$$\frac{T(Fn)}{T(n)} = F^3$$

$$\frac{T(n)}{n^3} = \frac{T(Fn)}{(Fn)^3}$$

O(2ⁿ) Algorithms

 $\begin{array}{l} T(n)=c \ 2^n \\ T(Fn)=c \ 2^{(Fn)} \end{array}$

$$\frac{T(Fn)}{T(n)} = 2^{(F-1)n}$$

$$\frac{T(n)}{2^n} = \frac{T(Fn)}{2^{(Fn)}}$$